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Reliable Key Distribution in Smart Micro-Grids

Heinrich Strauss^a, Anne VDM Kayem^a
and Stephen D Wolthusen^b

Presented by **Anesu Marufu^a**

^a *Department of Computer Science, University of Cape Town*

^b *NISLab, Faculty of Computer Science and Media Technology,
NTNU-Gjøvik*



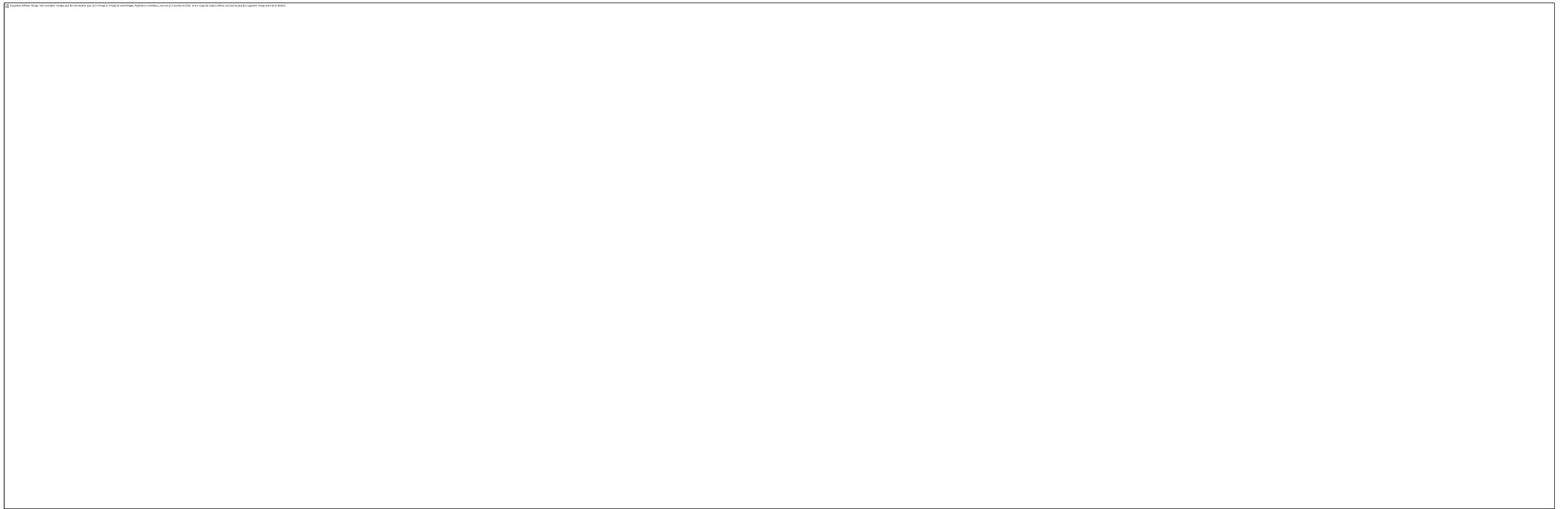
Outline

- Introduction
- Summary of Related Work
- Network Architecture
- Key Distribution Problem
- Correctness / Analysis Sketch
- Summary of Presentation
- References

Introduction

- Smart-Grids allow bi-directional information flow in a power network
- Micro-Grids enable use of self-generated power in isolated networks
- Smart Micro-grids combine the abovementioned concepts
 - Enable remote/mobile communities to manage a scarce resource
- How do we ensure attributable usage reports in a distributed, resource-constrained network?

Network – Physical Configuration



Problem

- User Trust in the network must exist for continued participation
 - If spurious reports are accepted as valid, this may be compromised
- Sensor Nodes are constrained
 - Power-resources are battery-backed (therefore limited)
 - Asymmetric cryptographic signatures are too expensive
- Network is not reliable
 - Based on Wireless links within and between locations
 - For constrained radio protocols, message sizes are small

Summary of Related Work

- Use Eschenauer-Gligor [1] algorithms
 - Probabilistic in the derivation of a shared-key
 - Expensive re-keying in the event of disclosure
- Symmetric Encryption
 - Pros: less intensive for constrained nodes
 - Cons: signing is still comparatively expensive
- Pre-Distribution and Delayed Authentication [2,3]
 - Nodes authenticate neighbours on critical path through the network

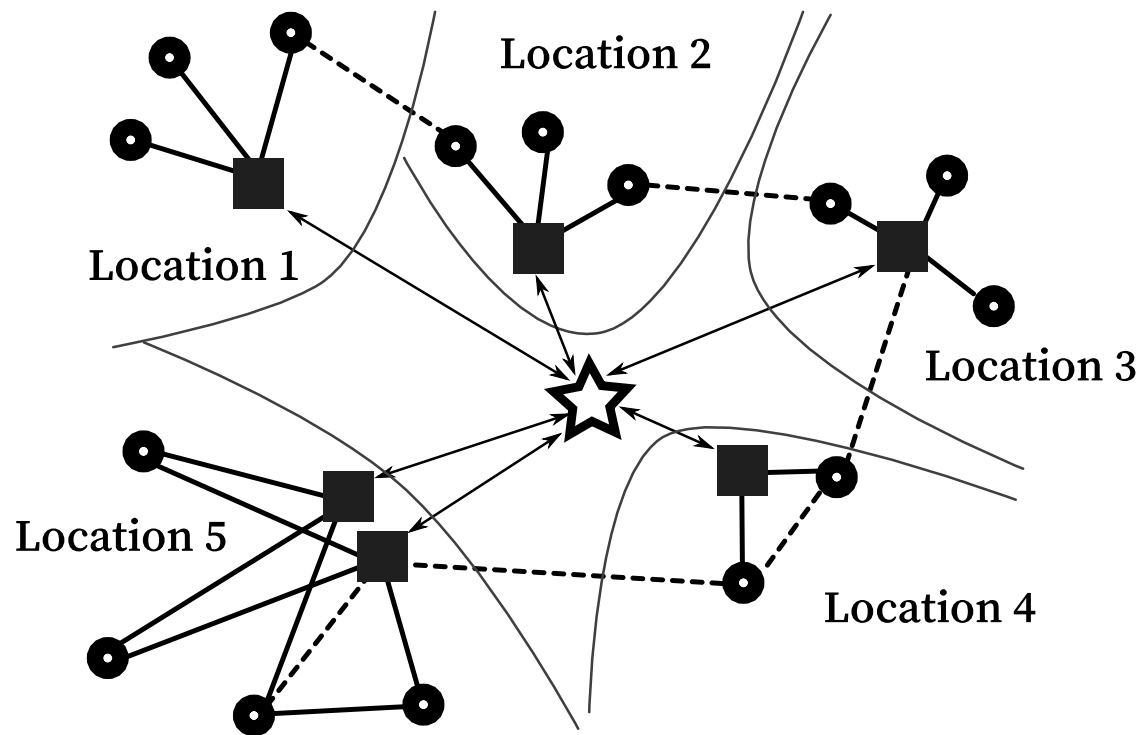
Requirements

- Unified neighbourhood-wide management network
 - For consolidation of aggregated usage reports
 - Security between devices ensured by Public-Key Cryptography (PKC)
- Per-location network for usage report collection
 - Usage report communication does not routinely extend beyond local network
 - Reports can be cached at sensor nodes in the event of short network outages
- Higher-powered Manager nodes in each location
 - Normally, a smart-device capable of asymmetric cryptographic operations [4]
 - Acts as gateway for nodes in each location

Key Distribution

- Key Pre-distribution
 - Central Global Pool of keys and key-tags
 - Partitioned into equal-sized location pools
 - Generated Centrally; distributed at node-addition time
- Key Revocation
 - Handled locally (per node)
- Key Validity
 - Keys are valid in a specific location only (determined by location pool)
 - Unless locally revoked, keys are considered valid
 - Manager nodes maintain a list of revoked keys per location

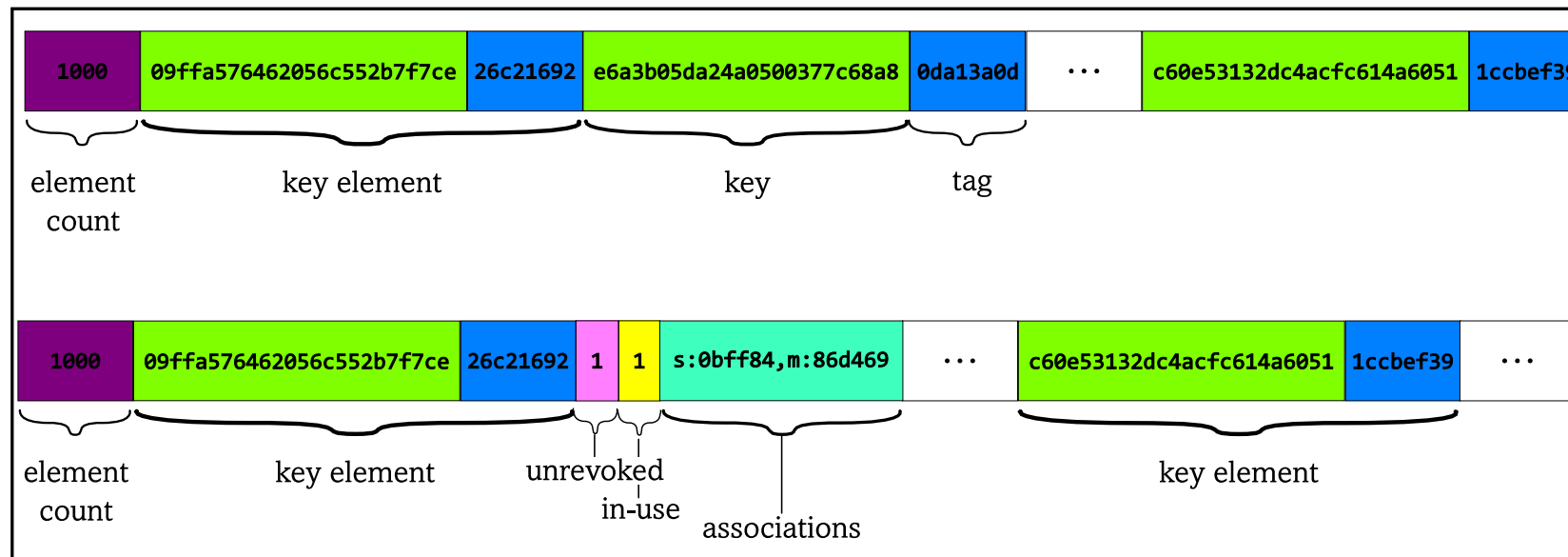
Network – Logical Configuration



- ☆ Utility Management Network
- Manager
- Sensor
- ↔ Neighbourhood Network
- shared key exists between nodes
- no shared key exists between nodes

Key Ring representation

- Representation dependent on location in network
 - At key-distribution, no metadata exists
 - On sensor/leaf nodes, usage data is included on key-ring



Correctness Analysis

- Network segmented into two classes: asymmetric and probabilistic symmetric
- Manager devices participate in both networks
 - PKI ensures security on the asymmetric network
 - Probabilistic EG-scheme ensures low probability of compromise in other network
- Barring node-compromise, keys are never disclosed in the clear over the transport medium
- If a node is known to be compromised, revocation can be initiated by the manager
 - List of compromised keys is distributed through the location
 - Damage restricted to single location (because of partitioning of key-pool)

Complexity Analysis

- The Key Distribution Center (KDC)
 - draws k random keys from a location pool ($O(k)$)
 - compares each key-tag to the list of key-tags issued to the n nodes in the location ($O(k \cdot n)$),
- This results in an algorithmic complexity of $O(n \times k^2) \approx O(n^3)$.
- Rekeying a k -key ring using a message size of μ , (key+tag)-length of ℓ
 - number of messages would be $n = \frac{k(\ell)}{\mu} + C$
- Online re-keying messages are very distinguishable from usage reporting traffic patterns

Conclusions – Summary and Future Work

- We can improve computational complexity by trading off heavy asymmetric encryption for symmetric encrypted between sensors and managers
- Managers act as gateways to node relaying messages to the central network
- Key Management Functions are maintained
 - Key Generation/Distribution
 - Key Revocation
 - Key Validity Checks

Selected References

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