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Domain Specific Stateful Filtering with Worst-Case Bandwidth

Maxime Puys, Marie-Laure Potet, and Jean-Louis Roch

*Verimag, University Grenoble Alpes/Grenoble-INP, Gières, France
Firstname.Name@imag.fr*



Industrial Systems (SCADA)



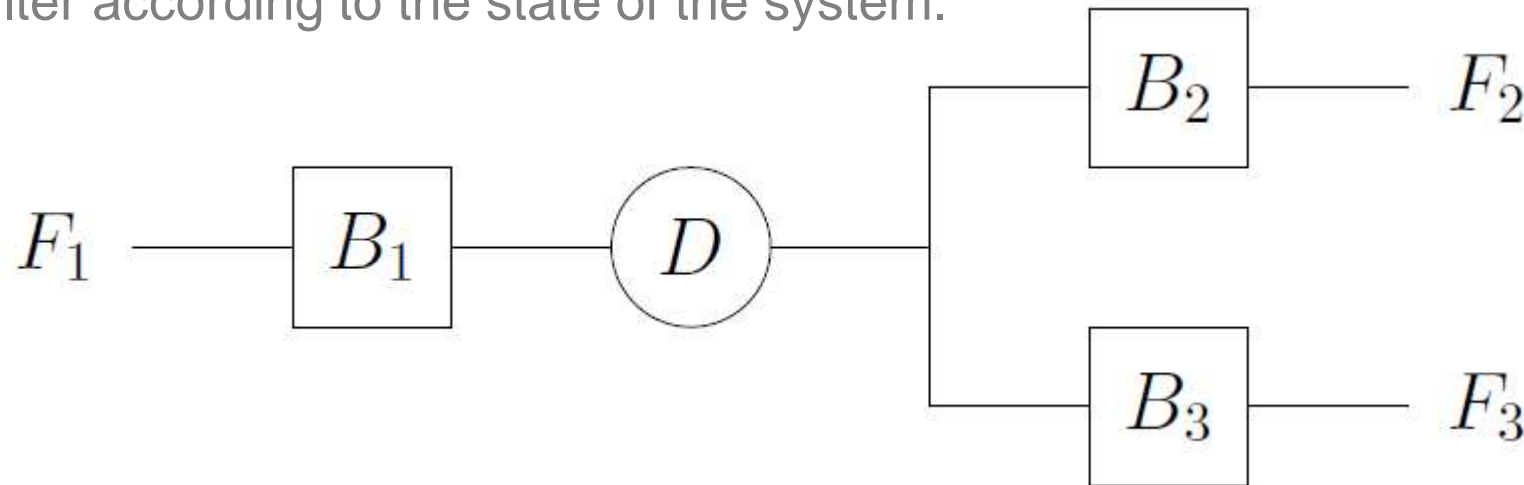
- Hot topic:
 - Increasing number of attacks showed in the media since Stuxnet.
 - Complex attack ending up in increasing speed of Iranian centrifuges to damage them.
 - Also attacked the process monitoring to trick operators.
 - Protection becoming a priority for government agencies.

Specificities of Industrial Systems

- Security objectives are different from traditional systems:
 - **Availability, integrity**, authentication and non-repudiation.
- Messages are READ/WRITE commands to PLCs.
 - Sometimes SUBSCRIPTIONS, RPCs or grouped commands.
 - Industrial protocols: MODBUS, OPC-UA.
- Attack examples: change the value of a WRITE request to change a temperature, change a READ response to mislead operators.

Motivating Example

- An electrical disconnector D :
 - Between three feeders F_1, F_2, F_3 .
 - Each separated by a breaker.
- No manipulation of D while current is passing through.
 - Either B_1 is open OR (B_2 AND B_3 are open).
 - Need to filter according to the state of the system.

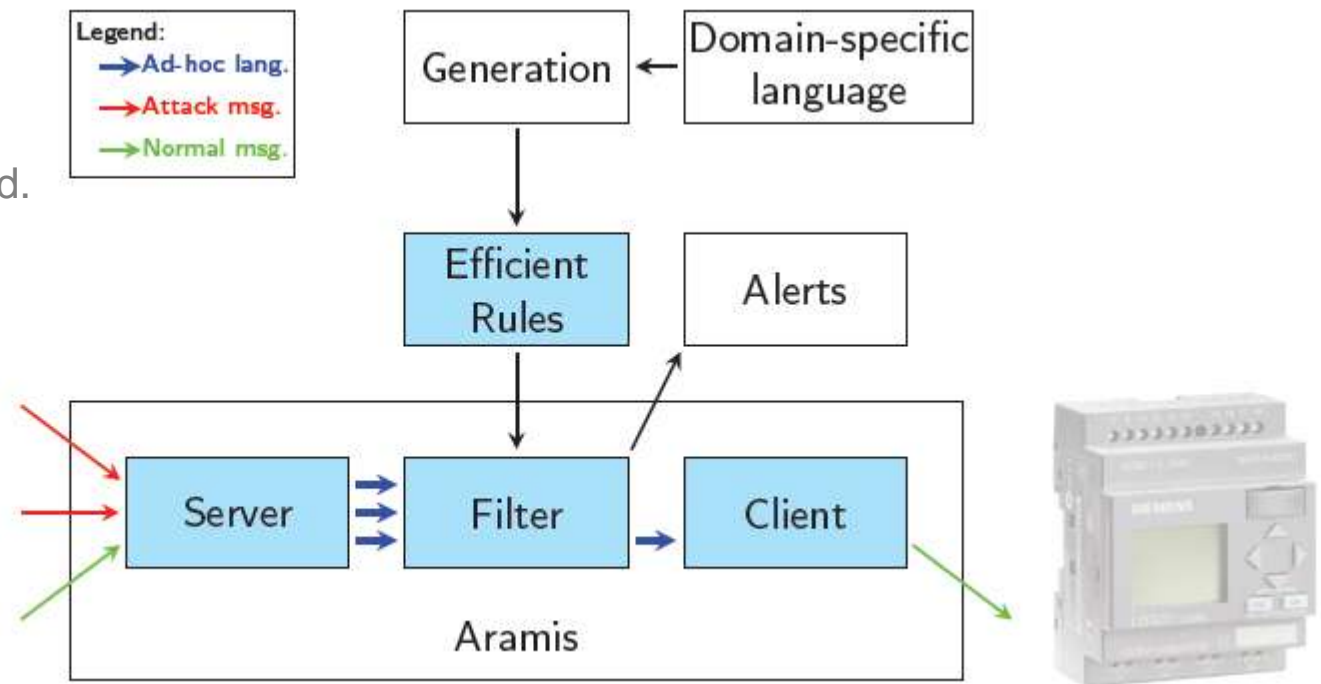


PIA ARAMIS Project

- Architecture Robuste pour les Automates et Matériels des Infrastructures Sensibles.
- Partners: ATOS World Grid, SecLab, CEA-leti, University of Grenoble Alpes.

Objectives:

- Build a transparent device to filter industrial flows.
- Communication protocols are stopped and re-created.
- Applicative filtering with domain specific constraints



Contributions

- Protocol-independent low-level language.
 - SCADA specific protocols (e.g.: MODBUS, OPC-UA).
- To describe a stateful type of domain specific filtering.
 - Automatically keeps track of the value of predetermined variables.
- Worst-case bandwidth and memory to meet real-time constraints.

Language 1/2

- Servers and Variables:

- Variables configured on the real server.

A MODBUS server.

Declare Server 1 Protocol **Modbus** Addr 10.0.0.1 Port 502

A MODBUS coil (read / write Boolean).

Declare Variable 1 Server 1 Type Boolean Addr **coils:0x10**

An OPC-UA server.

Declare Server 2 Protocol **OpcUa** Addr 10.0.0.2 Port 48010

An OPC-UA unsigned integer 5x10 matrix.

Declare Variable 2 Server 2 Type UInt32 Addr **numeric:5000** Dims 5 10

Language 2/2

- Domain specific safety rules to be verified:
 - Assert (= Block) or Warning (= Only logs).
 - Boolean predicates separated by AND and OR operators.

Variable 1 should never been set to its current value
(e.g.: opening a currently opened circuit breaker).

Declare Rule Variable 1 Assert NotEqual(NewVal, LocalVal[1])

Language 2/2

- Local state variables:
 - Local copy of server variables within the filter.
 - Updated on READ responses and WRITE responses (with requested values).
 - Only target scalar variables (use of Index keyword).

A local variable on a MODBUS coil.

Declare LocalVal 1 Variable 1

- Limited to local vision of the filter (at least the same as the client since we intercept communications).

Implementation

- Non-shared variables between clients:
 - Prevent isolation problem.
- When loaded into the filter, rules are indexed by (client, server) and then by variable:
 - Fast access to rules to verify.
- Embedded on the ARAMIS device.

Worst-Case Bandwidth

Both conditions and actions have to be processed in constant time:

- Conditions are $O(1)$ boolean predicates.
- Actions are : (i) Block or transmit the message, (ii) Log information, (iii) Update a local variable,

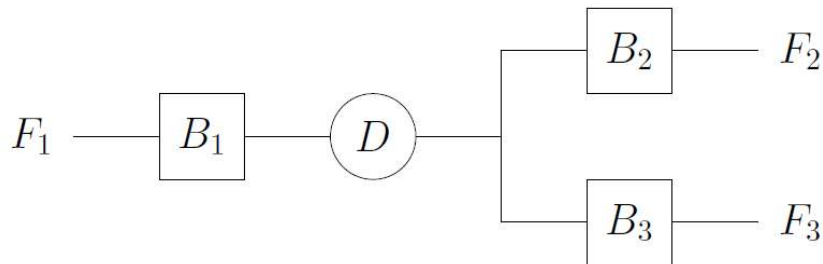
Thus processing one command only depends on the number of rules:

- For all predicates P , worst case processing time T of a message is $T = \sum \tau_i n_i$
- With τ_i the processing time of predicate P_i
- And n_i number of occurrences of predicate P_i

In practice, as only relevant rules are tested for a message, performances are better than upper-bound.

Motivating Example

D can be manipulated if and only if either B_1 is opened or if both B_2 and B_3 are open.



```
Declare Server 1 Protocol Modbus Addr 10.0.0.1 Port 502
```

```
Declare Variable 1 Server 1 Type Boolean Addr coils:0x11 # B1
```

```
Declare Variable 2 Server 1 Type Boolean Addr coils:0x12 # B2
```

```
Declare Variable 3 Server 1 Type Boolean Addr coils:0x13 # B3
```

```
Declare Variable 4 Server 1 Type Boolean Addr coils:0x14 # D
```

```
Declare LocalVal 1 Variable 1 # Local variable on B1
```

```
Declare LocalVal 2 Variable 2 # Local variable on B2
```

```
Declare LocalVal 3 Variable 3 # Local variable on B3
```

```
Declare Rule Variable 4 Assert \ # Rule on variable 4 = D
    Equal(LocalVal [1], False ) OR \ # Using local variable on B1, B2,
    B3
```

```
    Equal(LocalVal [2], False ) AND Equal(LocalVal[3], False )
```

Related Work

- Classical safety monitoring properties, e.g.: Roşu, 2012
 - Our language goes into an embedded device with high constraints.
 - Restrictions for filtering SCADA.
- Lot of IDS/IPS (e.g.: Snort or Bro):
 - Such IDS/IPS are unable to keep track of variable values.
- Chen et al., 2014:
 - Estimation through machine learning, MODBUS only
- Stergiopoulos et al., 2015:
 - Code review in industrial softwares to predict failures.

Conclusion

- Stateful filtering language:
 - Restriction of classic monitoring properties.
 - Keeps track of the value of predetermined variables.
- Bounded memory space and execution time:
 - Processing one command only depends on the number of rules.
 - Upper-bound allows to compare the worst-case path of the filter's configuration file with field constraints.
- Future work:
 - Handle more complex arithmetics : "Equal(2*NewVal+1, LocalVal[1]**2)".
 - Rules to avoid domain specific denial of service.
 - More precise upper-bound.

References

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Conclusion

Thanks for your attention!

`Maxime.Puys@imag.fr`